## Problem 1

First we formulate the two programs into KAT.

while b do {
 p;
 while c do q
}

can be expressed as

$$(bp(cq)^*\bar{c})^*\bar{b} \tag{1}$$

and

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if b then {

p;

while b \lor c do {

if c then q else p

}
```

can be expressed as

$$bp((b+c)(cq+\bar{c}p))^* \overline{b+c} + \bar{b}$$
<sup>(2)</sup>

Because  $\overline{b+c} = \overline{c} \, \overline{b}$  and

$$\begin{aligned} (b+c)(cq+\bar{c}p) &= bcq+b\,\bar{c}p+ccq+c\,\bar{c}p\\ &= bcq+b\,\bar{c}p+cq\\ &= (b+1)cq+b\,\bar{c}p\\ &= cq+\bar{c}bp, \end{aligned}$$

(2) can be transformed into

$$bp(cq + \bar{c}bp)^* \bar{c}\bar{b} + \bar{b} \tag{3}$$

Then we use the rule  $1 + pp^* = p^*$  on (1). We get

$$(bp(cq)^*\overline{c})^*\overline{b} = (1 + (bp(cq)^*\overline{c})(bp(cq)^*\overline{c})^*)\overline{b} = bp(cq)^*\overline{c}(bp(cq)^*\overline{c})^*\overline{b} + \overline{b}$$

$$(4)$$

This transformation correspond to loop unwinding operation, which turns the original program into

Now we can see that  $p(qp)^* = (pq)^* p$  may be applied on part in red in (4). We apply it and get

$$bp(cq)^*(\bar{c}\,bp(cq)^*)^*\bar{c}\,\bar{b}+\bar{b}$$

We see that we can apply this rule once more and get

$$bp((cq)^*\overline{c}\,bp)^*(cq)^*\overline{c}\,\overline{b}+\overline{b}$$

Now we can see the rule  $(p^*q)^*p^* = (p+q)^*$  can be applied to the red part with substitutions p = cq and  $q = \overline{c} bp$ . Apply it and we get

$$bp(cq+\overline{c}bp)^*\overline{c}\overline{b}+\overline{b}$$

which equals (3). This proves the programs are equal.